Homework 3

Algorithms for Elementary Algebraic Geometry Math 191, Fall Quarter 2007

Solutions.

1. Here is an an algorithm to solve the ideal membership problem in k[x]: Let polynomials f and f_1, \ldots, f_s in k[x] be given. Compute a greatest common divisor g of f_1, \ldots, f_s using the Euclidean Algorithm. If g divides f, then output " $f \in \langle f_1, \ldots, f_s \rangle$ "; otherwise output " $f \notin \langle f_1, \ldots, f_s \rangle$." (This is justified by $\langle g \rangle = \langle f_1, \ldots, f_s \rangle$.) Now we use this procedure and Maple to decide whether in $\mathbb{Q}[x]$ we have

$$x^{2} - 4 \in \langle x^{3} + x^{2} - 4x - 4, x^{3} - x^{2} - 4x + 4, x^{3} - 2x^{2} - x + 2 \rangle.$$

First we compute

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> gcd(x^3+x^2-4*x-4, x^3-x^2-4*x+4);

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x - 4

and then
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Hence

$$g = \text{GCD}(x^3 + x^2 - 4x - 4, x^3 - x^2 - 4x + 4, x^3 - 2x^2 - x + 2) = x - 2,$$

and since g divides $f = x^2 - 4 = (x + 2)(x - 2)$, we see that f lies in the ideal in question.

2. Our trusted companion Maple gives us:

Hence we are tempted to conjecture that in general,

$$GCD(x^m - 1, x^n - 1) = x^d - 1$$

where d > 0 is the greatest common divisor of the integers m and n.

3. Let $f \in \mathbb{C}[x], f \neq 0$.

(a) The proof is by induction on d = deg(f). If deg(f) = 0, then f is a constant, so f = c is a factorization of the desired form. Suppose that d > 0, and the claim is true for all polynomials of degree less than d. Assume that f has degree d. By the Fundamental Theorem of Algebra, when f is non-constant, f has a zero, say a, in C: f(a) = 0. Now the Division Algorithm yields that

$$f(x) = g(x)(x-a) + r(x)$$

where r = 0 or $\deg(r) < \deg(x - a) = 1$, so r is a constant. Thus 0 = f(a) = g(a)(a - a) + r, which implies that f = (x - a)g. Since

$$d = \deg(f) = \deg(g) + \deg(x - a) = \deg(g) + 1 > \deg(g),$$

by inductive hypothesis applied to g, there is a factorization of g in the form

$$g = c(x - a_1)^{r_1} \cdots (x - a_m)^{r_m}$$

where $c \in \mathbb{C}$ is nonzero, and a_1, \ldots, a_m are pairwise distinct. So

$$f = c(x-a)(x-a_1)^{r_1}\cdots(x-a_m)^{r_m}$$

If a, a_1, \ldots, a_m are pairwise distinct, this is a factorization of f in the required form. If $a = a_i$ for some i then

$$f = c(x - a_1)^{r_1} \cdots (x - a_i)^{r_i + 1} \cdots (x - a_m)^{r_m}$$

is the desired factorization.

[Note that there was still another small inaccuracy in how the problem was formulated: instead of r_1, \ldots, r_m non-negative, they should be required to be positive!]

- (b) Clearly if $a \in \{a_1, a_2, \ldots, a_m\}$, then f(a) = 0. If $a \notin \{a_1, \ldots, a_m\}$ then $f(a) = c(a a_1)^{r_1} \cdots (a a_m)^{r_m}$ is the product of nonzero elements of \mathbb{C} , so is nonzero. Thus $V(f) = \{a_1, a_2, \ldots, a_m\}$.
- (c) Since $f_{\text{red}}(a_i) = 0$ for $1 \leq i \leq m$, we have $f_{\text{red}} \in I(V(f))$. For the reverse inclusion, let $g \in I(V(f))$. Then $g(a_i) = 0$ for $1 \leq i \leq m$. We claim that g is a multiple of f_{red} . The proof is by induction on the size m of V(f). If m = 1, then $f_{\text{red}} = x a_1$, and the Division Algorithm implies that that

$$g = q(x - a_1) + r,$$

where r is a constant; the fact that $g(a_1) = 0$ means that r = 0, so g is a multiple of $x - a_1$. Now suppose that the claim is true if V(f) < m. The polynomial

$$f_1 := (x - a_1) \cdots (x - a_{m-1})$$

satisfies $V(f_1) = \{a_1, \ldots, a_{m-1}\}$; hence we have $I(V(f_1)) = \langle f_1 \rangle$ since $(f_1)_{\text{red}} = f_1$. Since $g(a_i) = 0$ for $1 \leq i \leq m-1$, we know that $g = f_1 h$ where $h \in \mathbb{C}[x]$. Since $g(a_m) = 0$ but $f_1(a_m) = (a_l - a_m) \cdots (a_m - a_{m-1}) \neq 0$, we must have $h(a_m) = 0$, and so by the base case $h = (x - a_m)p$ for some $p \in \mathbb{C}[x]$. Thus

$$g = f_1 h = (x - a_1) \dots (x - a_{m-1})(x - a_m)p$$

is a multiple of f_{red} . This shows that $I(V(f)) \subseteq \langle f_{\text{red}} \rangle$, and so the two ideals are equal.

(d) One first checks by computation that the operation $p \mapsto p'$ satisfies the usual properties of the derivative: for all $p, q \in \mathbb{C}[x]$ we have

(p+q)' = p'+q', (pq)' = p'q + pq' (Product Rule).

This implies that $(p^n)' = np^{n-1}p'$ for every positive integer n and $p \in \mathbb{C}[x]$. (I omit some details here.) Now applying the product rule to

$$f = c(x - a_1)^{r_1} \cdots (x - a_m)^{r_m}$$

we obtain

$$f' = cr_1(x - a_1)^{r_1 - 1}(x - a_2)^{r_2} \cdots (x - a_m)^{r_m} + cr_2(x - a_1)^{r_1}(x - a_2)^{r_2 - 1} \cdots (x - a_m)^{r_m} + \dots + cr_m(x - a_1)^{r_1}(x - a_2)^{r_2} \cdots (x - a_m)^{r_m - 1} = c(x - a_1)^{r_1} \cdots (x - a_m)^{r_m} \left(\frac{r_1}{x - a_1} + \dots + \frac{r_m}{x - a_m}\right) = c(x - a_1)^{r_1 - 1} \cdots (x - a_m)^{r_m - 1} H(x)$$

where

$$H = (x - a_1) \cdots (x - a_m) \left(\frac{r_1}{x - a_1} + \cdots + \frac{r_m}{x - a_m} \right).$$

Then H is a polynomial (why?) such that $H(a_i) \neq 0$ for any i. This implies that

$$GCD(f, f') = (x - a_1)^{r_1 - 1} \cdots (x - a_m)^{r_m - 1}.$$

- (e) This follows from (a) and (d). [Note that strictly speaking, this equation is only true "up to multiplication by c." Perhaps I should have assumed from the beginning that f is monic, since then c = 1.]
- (f) Using Maple, we compute the formal derivative of

$$f = x^{11} - x^{10} + 2x^8 - 4x^7 + 3x^5 - 3x^4 + x^3 + 3x^2 - x - 1$$

as follows:

Hence by (c):

$$I(V(x^{11}-x^{10}+2x^8-4x^7+3x^5-3x^4+x^3+3x^2-x-1)) = \langle x^5+x^2-x-1\rangle.$$